



Module Based Data Management in Secured Procedure

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Abstract:

The present pattern for implementing a philosophy based data administration framework is to gain by efforts made to outline a preexisting well-secured Data Management System. The technology adds up to concentrating from the reference DMS a bit of mapping pertinent to the new application needs a module, possibly customizing it with additional obligations to the application under development, and after that dealing with information set utilizing the ensuing mapping. The current meanings of modules and we present novel properties of robustness that give intends to checking effectively that a vigorous module-based DMS develops securely both the pattern and the information of the reference DMS. The improved methodology helps in circulated module-based Data, where replying inquiries consolidates information of a few modules associated with possibly a few reference DMS. Modularized DMS in appropriated way needs to be classified in numerical path in which inquiry noting could likewise be integrated.

1. INTRODUCTION:

In many application domains comprehensive schemas resulting from collaborative initiatives are made available. Such well-established schemas are often associated with reliable data that have been carefully collected, cleansed, and verified, thus providing reference ontology-based data management systems in different application domains.

A good practice is therefore to build on the efforts made to design reference DMSs whenever we have to develop our own DMS with specific needs. A way to do this is to extract from the reference DMS the piece of schema relevant to our application needs, possibly to personalize it with extra-constraints our application under construction, and then to manage our own dataset using the resulting schema. Recent work in description logics provides different solutions to achieve such a re-use of a reference ontology-based DMS. Indeed, modern

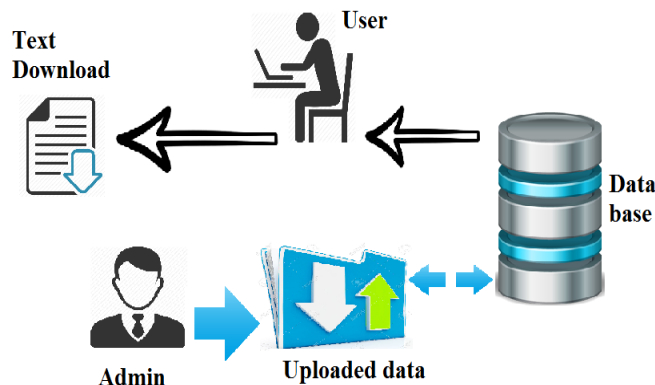


ontological languages like the W3C recommendations OWL, and OWL2 are actually XML based syntactic variants of well-known DLs.

All those solutions consist in extracting a module from an existing ontological schema such that all the constraints concerning the relations of interest for the application under construction are captured in the module. Existing definitions of modules in the literature basically resort to the notion of deductive conservative extension of a schema or of uniform interplant of a schema forgetting about non-interesting relations of a schema.

Formalizes those two notions for schemas written in DLs and discusses their connection. Up to now, conservative extension has been considered for defining a module as a subset of a schema. In contrast, forgetting has been considered for defining a module as only logically implied by a schema. Both kinds of modules have been investigated in various DL.

2. Architecture:



The architecture mainly describes about the user accessed text data to database. The uploaded data will be authorized by the admin. Admin providing security and database.

3. Modules Description:

3.1 User Module:

Users are provided authentication and security to access the files which is presented in the ontology system. Before accessing or searching the details user should have the account in that otherwise they should register first.



3.2 Global Answer Illustration:

Suppose now that our DMS can answer conjunctive queries asking for the journal papers written by Alon Y. Halevy. In some situation, it is interesting to provide answers from our DMS together with the reference one, called global answers, typically when our own DMS provides no or too few answers. To do so, we extend the notion of module to robustness to query answering, so that global query answering can be performed on demand. We ensure that the module captures the knowledge in the reference schema that is required to answer any query built upon the relations of interest. Then, at global query answering time, this knowledge is used to identify the relevant data for a given query within the distributed dataset consisting of the dataset of the module-based DMS plus that of the reference DMS.

3.3 Reducing Data Storage Illustration:

Computing edit distance exactly is a costly operation. Several techniques have been proposed for identifying candidate strings within a small edit distance from a query string fast. All of them are based on q-grams and a q-gram counting argument. For a string s , its q-grams are produced by sliding a window of length q over the characters of s . To deal with the special case at the beginning and the end of s , that have fewer than q characters, one may introduce special characters, such as “#” and “\$”, which are not in S . This helps conceptually extends by prefixing it with $q - 1$ occurrences of “#” and suffixing it with $q - 1$ occurrences of “\$”. Hence, each q-gram for the string s has exactly q characters.

3.4 Module-Based Data Management:

The main idea underlying the notion of module of a Text is to capture some constraints of the Text box, including all the constraints built upon a given signature, denoted the signature of interest. Our definition of module extends and encompasses the existing definitions. In contrast with we do not impose modules of a Tbox to be subsets of it. For a module to capture some constraints of the Tbox, it is indeed sufficient to impose that it is logically entailed by the Tbox. In contrast with, we do not impose the signature of modules to be restricted to the signature of interest. In fact, as we have shown through the illustrative example, the robustness properties may enforce the signature of modules to contain additional relations that are not relations of interest but that are logically related to them.



Think India Journal

ISSN: 0971-1260 Vol-22, Special Issue-21

National Conference on

Recent Advances in Commerce, Management and Computer Science (NRCACMC-2020) organised

by

Department of Commerce, VEL TECH Ranga Sanku Arts College,
Avadi, Chennai-62

Held on 4th January 2020



4. Frame work



Fig: 4.1 Home page

It display user login and registration .Choose the required that is already registered then go to user login. Otherwise register newly then only user login successfully.



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Fig: 4.2 login page

User have registered then goto login page. It requires username, password. click on login button then only login page open.

ID	Name	Email	Password	DOB	Gender	Address
1	male	m@gmail.com	12345	3-Jul-2013	Null	Chennai, mylapore
2	m	m@gmail.com	1	3-Jul-2013	Male	Chennai
3	kannan	k@gmail.com	12345	15-Jul-1985	Null	Chennai
4	raja	r@gmail.com	12345	1-Jul-1985	Male	Bangalore
5	rakesh	r@gmail.com	revathi	18-08-1993	Male	London

Fig: 4.3 Admin page



It shows admin page containing the data present in database.

Fig: 4.4 File upload page

It shows how to upload the file with submit and reset option , for uploading select the file and submit .

5. CONCLUSION:

The modules introduced in this generalize both the modules obtained by extracting a subset of selected relations. In addition, in contrast with existing work, we have considered the problem of safe personalization of modules built from an existing reference DMS. This raises new issues to check easily that a module-based DMS evolves independently but coherently with respect to the reference DMS from which it has been built. We have introduced two notions of module robustness that make possible to build locally the relevant queries to ask to the reference database in order to check global consistency, and to obtain global answers for local queries. We have provided polynomial time algorithms that extract minimal and robust modules from a reference ontological schema expressed as a DL-lite Tbox. while keeping data consistency and query answering reducible to standard database queries In contrast with the recent work on extracting modules from DL-lite ontological schema focus on the DL-liteA fragment for which consistency checking and query answering are FOL-reducible. This is crucial when ontologies



are used as schemas over large datasets stored and queried as relational databases. Contrarily to recent works in distributed databases, data replication can be avoided while guaranteeing global consistency. Our approach is a good trade-off between the No SQL approaches and the SQL approaches for managing distributed data stores. We also plan to apply our algorithms to the real usecase of the My Corporis Fabrica DMS, mentioned in the introduction, which has been developed manually as a personalization of the Foundational Model of Anatomy DMS. Finally, we plan to extend our approach to distributed module-based DMSs, where answering queries combines knowledge of several modules associated with possibly several reference DMSs.

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